## Étiquetage des régions dans les arrangements d'hyperplans

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(1) Preliminaries

- Hyperplane arrangements
- Zaslavsky's formulas
- Inequality based approaches
(2) Inequalities for deformed graphical arrangements
- The general setup
- Sparse deformations
- Separated deformations

Outline

Hyperplane arrangements
Zaslavsky's formulas
Inequality based approaches

## Hyperplane arrangements

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A hyperplane arrangement $\mathcal{A}$ is a finite collection of hyperplanes in a d-dimensional real vector space, which partition the space into regions.

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1 bounded and 6 unbounded regions

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## Deformations of the braid arrangement

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The braid arrangement (Coxeter arrangement of type $A_{n-1}$ ) is the collection of hyperplanes $\left\{x_{i}-x_{j}=0: 1 \leq i<j \leq n\right\}$ in $V_{n-1}$, the subspace of $\mathbb{R}^{n}$, given by $x_{1}+x_{2}+\cdots+x_{n}=0$.

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x_{i}-x_{j}=a_{i j}^{(1)}, a_{i j}^{(2)}, \ldots, a_{i j}^{\left(n_{i j}\right)}
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The truncated affine arrangements $\mathcal{A}_{n-1}^{a, b}$ (where $a+b \geq 2$ ) contain the hyperplanes are $x_{i}-x_{j}=1-a, 2-a, \ldots, b-1$ for $1 \leq i<j \leq n$.

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Outline

## The characteristic polynomial

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To count the regions, we may use Zaslavsky's formulas ("inclusion-exclusion") or solve systems of linear inequalities directly.

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and compute the characteristic polynomial
$\chi(\mathcal{A}, q)=\sum_{x \in L_{\mathcal{A}}} \mu(\widehat{0}, x) q^{\operatorname{dim}(x)}=1-3 q+3 q^{2}$.

Outline

Hyperplane arrangements Zaslavsky's formulas
Inequality based approaches

## Zaslavsky's formulas

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The numbers $r(\mathcal{A})$ and $b(\mathcal{A})$ of all, respectively bounded regions are given by

$$
r(\mathcal{A})=(-1)^{d} \chi(\mathcal{A},-1) \quad \text { and } \quad b(\mathcal{A})=(-1)^{\text {rk }\left(L_{\mathcal{A}}\right)} \chi(\mathcal{A}, 1) .
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In our example

$$
r(\mathcal{A})=(-1)^{2}\left(1-3 \cdot(-1)+3 \cdot(-1)^{2}\right)=7
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and

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b(\mathcal{A})=(-1)^{2}(1-3+3)=1
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Related approaches: finite field method (case of integer coefficients), Whitney's formula and the gain graph method (deformations of graphical arrangements).

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x_{2}-x_{3}>1 \\
x_{1}-x_{3}>1
\end{array} \quad\left(\begin{array}{l}
x_{1}-x_{2}>1 \\
x_{2}-x_{3}<1 \\
x_{1}-x_{3}>1 \\
(1,0,-1)
\end{array} \quad / \begin{array}{l}
x_{1}-x_{2}>1 \\
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x_{1}-x_{3}<1 \\
(2 / 3,-1 / 3,-1 / 3)
\end{array}\right. \\
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\end{array} \\
& x_{1}-x_{3}>1 \\
& x_{2}-x_{3}<1 \\
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\end{aligned}
$$

One possibility is missing:

## Regions defined by sets of inequalities

$$
x_{1}-x_{2}>1 \text { and } x_{2}-x_{3}>1 \text { imply } x_{1}-x_{3}>1
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## Examples of the inequality based approach

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The general setup Sparse deformations
Separated deformations

## Two key lemmas

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The following variant of the Farkas Lemma was also used by Hopkins and Perkinson:

## Lemma (Carver)

The system of inequalities $A x<b$ has no solution if and only if there is a nonzero real $m \times 1$ row vector $y$ satisfying $y \geq 0, y A=0$ and $y b \leq 0$.

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We will apply the flow decomposition theorem to circulations:

## Theorem (Gallai)

Every not identically zero circulation $f$ can be written as a positive linear combination of directed cycles. Moreover, a directed edge e appears in at least one of these cycles if and only if $f(e)>0$.

## Weighted digraphical polytopes

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A weighted digraphical polytope is the solution set of a system of inequalities

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m_{i j}<x_{i}-x_{j}<M_{i j}, \quad 1 \leq i<j \leq n
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in $V_{n-1}$. (We allow $m_{i j}=-\infty$ and $M_{i j}=\infty$.)
We create an associated weighted digraph: For each $i<j$, if $m_{i j}>-\infty$, we create directed edge $i \rightarrow j$ with weight $m_{i j}$ and if $M_{i j}<\infty$ we also create a directed edge $i \leftarrow j$ with weight $-M_{i j}$.

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## Proof.

(Sketch) By Carver's variant of the Farkas Lemma the polytope is empty if and only if there is an " $m$-ascending circulation". By the Flow Decomposition Theorem every $m$-ascending circulation contains an $m$-ascending cycle.

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## Corollary

If we think of the weight $w(e)$ as money we gain when we walk along e then the system of inequalities has a nonempty solution set if and only if we lose money along any closed walk.

Outline

## Semiacyclic tournaments

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If all arrows go from $V_{2}$ to $V_{1}$ then $\left(x_{1}, \ldots, x_{n}\right)$ may be replaced with $\left(x_{1}^{\prime}, \ldots, x_{n}^{\prime}\right)$ where

$$
x_{v}^{\prime}= \begin{cases}x_{v}+\frac{t}{\left|V_{1}\right|} & \text { if } v \in V_{1} \\ x_{v}-\frac{t}{\left|V_{2}\right|} & \text { if } v \in V_{2}\end{cases}
$$

## Bounded regions

## Theorem

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## Example

Each region of the Linial arrangement is described by a set of inequalities $\left\{m_{i j}<x_{i}-x_{j}<M_{i j}: 1 \leq i<j \leq n\right\}$, each inequality is either $-\infty<x_{i}-x_{j}<1$ or $1<x_{i}-x_{j}<\infty$. The associated weighted digraph is a tournament, it contains no m-ascending cycle if and only if it is semiacyclic. Bounded regions correspond to strongly connected semiacyclic tournaments.

## Exponential arrangements

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Let $\mathcal{A}=\left(\mathcal{A}_{1}, \mathcal{A}_{2}, \ldots\right)$ be a sequence of deformations of the braid arrangement, such that each $\mathcal{A}_{n}$ is a hyperplane arrangement in $\mathbb{R}^{n}$. For each $S \subseteq\{1,2, \ldots\}$ we define $\mathcal{A}_{n}^{S}$ as the subcollection of hyperplanes $x_{i}-x_{j}=c$ of $\mathcal{A}_{n}$ satisfying $\{i, j\} \subseteq S$. $\mathcal{A}$ is exponential if $r\left(\mathcal{A}_{n}^{S}\right)$ depends only on $k=|S|$ and it is the number $r\left(\mathcal{A}_{k}\right)$ of regions of $\mathcal{A}_{k}$.

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$$
B_{\mathcal{A}}(t)=1-\frac{1}{R_{\mathcal{A}}(t)}
$$

## Exponential arrangements (cont'd)

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$$
r\left(\mathcal{A}_{n}\right)=\sum_{k=1}^{n} \sum_{\substack{n_{1}+\cdots+n_{k}=n \\ n_{1}, \ldots, n_{k}>0}}\binom{n}{n_{1}, n_{2}, \ldots, n_{k}} \prod_{i=1}^{k} b\left(\mathcal{A}_{n_{i}}\right) \quad \text { for all } n \geq 1 .
$$

## Posets of gains

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## Definition

Given a valid $m$-acyclic weighted digraph $D$ on $\{1,2, \ldots, n\}$, we define $i<_{D} j$ if there is a directed path $i=i_{0} \rightarrow i_{1} \rightarrow \cdots \rightarrow i_{k}=j$ such that the weight of each directed edge $i_{s} \rightarrow i_{s+1}$ is nonnegative. We call the set $\{1,2, \ldots, n\}$, ordered by ${<_{D}}_{D}$ the poset of gains induced by $D$.

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The relation $i<_{D} j$ is a partial order because of the $m$-acyclic property.

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## Example

The posets of gains of the Linial arrangement are the sleek posets.

## Sparse deformations

## Sparse deformations

## Definition

a deformation of the braid arrangement, is sparse if $1 \leq n_{i, j} \leq 2$ holds for all $i<j$, and the signs of the numbers $a_{i, j}^{(k)}$ satisfy the following for all $i<j$ :
(1) $a_{i, j}^{(1)}>0$ holds, whenever $n_{i, j}=1$,
(2) $a_{i, j}^{(1)}<0<a_{i, j}^{(2)}$ holds, whenever $n_{i, j}=2$.

We call $\mathcal{A}$ an interval order arrangement if $n_{i, j}=2$ holds for all $i<j$.

## Sparse deformations

## Proposition

Consider a sparse deformation of the braid arrangement and any valid m-acyclic weighted digraph $D$ associated to it. In the induced poset of gains, $i<_{D} j$ holds exactly when there is a single directed edge $i \rightarrow j$ of positive weight. For any pair $\{i, j\}$ of incomparable vertices satisfying $i<j$, the edge $j \rightarrow i$ is always present, and any edge between $i$ and $j$ has negative weight.

## Sparse deformations

## Theorem

Let $D$ be a valid m-acyclic weighted digraph associated to a sparse deformation of the braid arrangement in $V_{n-1}$. If $D$ is strongly connected then the incomparability graph of the induced poset of gains is connected. The converse is also true when $n_{i, j}=2$ holds for all $1 \leq i<j \leq n$.

## Sparse deformations

## Example

Consider the Linial arrangement and the semiacyclic tournament $D$ containing a directed edge $i \leftarrow j$ of weight -1 for each $i<j$. This is a valid $m$-acyclic weighted digraph, it is in fact acyclic. The induced poset of gains is an antichain, the incomparability graph is the complete graph, it is connected. However, $D$ is not strongly connected.

## a-generalized Linial arrangements

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## Definition

Let $\underline{a}=\left(a_{1}, a_{2}, \ldots, a_{n}\right) \in \mathbb{R}_{\geq 0}{ }^{n}$. The $\underline{a}$-generalized Linial arrangement is

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x_{i}-x_{j}=a_{i} \quad \text { for } 1 \leq i<j \leq n \text { in } V_{n-1} .
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If $D$ is a valid m-acyclic weighted digraph associated to an a-generalized Linial arrangement, then $D$ contains no alternating cycle.

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## Proposition

If $D$ is a valid m-acyclic weighted digraph associated to an a-generalized Linial arrangement, then $D$ contains no alternating cycle.

Alternation acyclic tournaments label the regions of the homogenized Linial arrangement $\left\{x_{i}-x_{j}=y_{j}: 1 \leq i<j \leq n\right\}$.

Outline

## Separated deformations

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## Definition

We call a deformation of the braid arrangement $\mathcal{A}$ separated if 0 belongs to the set $\left\{a_{i j}^{(1)}, a_{i j}^{(2)}, \ldots, a_{i j}^{\left(n_{i j}\right)}\right\}$ for each $1 \leq i<j \leq n$.

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## Corollary

For a separated deformation of the braid arrangement, the induced poset of gains associated to any valid m-acyclic weighted digraph is a totally ordered set.

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Equivalently, each region is included in a region $x_{\sigma(1)}>x_{\sigma(2)}>\cdots>x_{\sigma(n)}$ of the braid arrangement.

Outline

## A structure theorem

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## Theorem

Let $\mathcal{R}$ be a region of a separated deformation of the braid arrangement and let $\sigma(1) \sigma(2) \cdots \sigma(n)$ be its total order of gains.
Then there is a unique decomposition
$\sigma=\left(\sigma\left(i_{0}\right) \cdots \sigma\left(i_{1}\right)\right) \cdot\left(\sigma\left(i_{1}+1\right) \cdots \sigma\left(i_{2}\right)\right) \cdots\left(\sigma\left(i_{k-1}+1\right) \cdots \sigma\left(i_{k}\right)\right)$
satisfying
(1) For each $j=-1,0, \ldots, k-1$,
$\mathcal{R} \cap \operatorname{span}\left(e_{\sigma\left(i_{j}+1\right)}, e_{\sigma\left(i_{j}+2\right)}, \ldots, e_{\sigma\left(i_{j+1}\right)}\right)$ is bounded.
(2) If $S \subseteq\{1,2, \ldots, n\}$ contains indices $j_{1}$ and $j_{2}$ such that $\sigma\left(j_{1}\right)$ and $\sigma\left(j_{2}\right)$ belong to different subwords in the above decomposition then $\mathcal{R} \cap \operatorname{span}\left(\left(e_{\sigma(j)}: j \in S\right)\right.$ is unbounded.

## Gain functions

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## Definition

For each $i \in\{1,2, \ldots, n\}$ we define the gain function $g(\sigma(i))$ as the maximum weight of a directed path beginning at $\sigma(1)$ and ending at $\sigma(i)$. In particular, we set $g(\sigma(1))=0$. Here $\sigma$ is the total order of gains.

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## Lemma

Every gain function has the weakly increasing property

$$
g(\sigma(1)) \leq g(\sigma(2)) \leq \cdots \leq g(\sigma(n))
$$

## Gain functions

## Definition

We call a deformation $\mathcal{A}$ of the braid arrangement integral if all the numbers $a_{i, j}^{k}$ appearing in in its definition are integers. We say that $\mathcal{A}$ satisfies the weak triangle inequality if for all triplets $(i, j, k)$, the inequalities $w(i, j) \geq 0$ and $w(j, k) \geq 0$ imply

$$
w(i, k) \leq w(i, j)+w(j, k)+1
$$

in any valid $m$-acyclic associated weighted digraph.

## Gain functions

## Theorem

Let $\mathcal{A}$ be a separated integral deformation of the braid arrangement satisfying the weak triangle inequality, and let $D$ be an associated m-acyclic weighted digraph. Let $\sigma$ be the total order of gains associated to $D$ and let $g$ be the gain function. Then, for each $i>1$ there is a directed path from $\sigma(1)$ to $\sigma(i)$ such that all weights in the path are nonnegative and the total weight of the edges in the path is $g(\sigma(i))-g(\sigma(1))$.

Outline

## Contiguous integral deformations

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## Definition

An integral deformation of the braid arrangement in $V_{n-1}$ contiguous if, for every $i<j$, the set $\left\{a_{i, j}^{(1)}, a_{i, j}^{(2)}, \ldots, a_{i, j}^{\left(n_{i, j}\right)}\right\}$ is a contiguous set $[\alpha(i, j), \beta(i, j)]=\{\alpha(i, j), \alpha(i, j)+1, \ldots, \beta(i, j)\}$ of integers.

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Since $x_{i}-x_{j}=c \Leftrightarrow x_{j}-x_{i}=-c$, we may set

$$
\alpha(j, i)=-\beta(i, j) \quad \text { and } \quad \beta(j, i)=-\alpha(i, j) \quad \text { for } 1 \leq i<j \leq n
$$

## Minimal obstructions

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## Theorem

If $\beta(i, k) \leq \beta(i, j)+\beta(j, k)+1$ holds for all $\{i, j, k\}$. then any valid associated weighted digraph is $m$-acyclic if and only if it contains no m-ascending cycle of length at most four.

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## Theorem

If the truncated affine arrangement $\mathcal{A}_{n-1}^{a, b}$ satisfies $a, b \geq 0$, then a valid associated weighted digraph is $m$-acyclic if and only if it contains no m-ascending cycle of length at most four.

## Minimal obstructions

There is a minimal $m$-ascending cycle of length 5 in $\mathcal{A}_{n-1}^{-1,3}$ for $n \geq 5$.


Outline

## The Pak-Stanley labeling

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We define the Pak-Stanley label $(f(1), \ldots, f(n))$ of a region as

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f(i)=\sum_{i<_{\sigma-1} j} w(i, j)+\mid\left\{(i, j): i<_{\sigma^{-1}} j \text { and } i>j\right\} \mid .
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The sum $\sum_{i<_{\sigma^{-1}} j} w(i, j)$ is the number of separations, and $\mid\left\{(i, j): i<_{\sigma^{-1}} j\right.$ and $\left.i>j\right\} \mid$ is the number of inversions.

## The Pak-Stanley labeling

## Lemma (Stanley)

Given $i<_{\sigma^{-1}} j$, if $i>j$ or $w(i, j)>0$ holds then we have $f(i)>f(j)$.

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The labels of the regions of the extended Shi arrangement are the a-parking functions of length n, each occurring exactly once.

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The labels of the regions of the extended Shi arrangement are the a-parking functions of length $n$, each occurring exactly once.

Given an a-parking function $(f(1), \ldots, f(n))$, we insert the labels $i$ into $\sigma$ one by one and show the uniqueness of the place and of the function values $w(i, j)$ one step at a time. (Still "tedious", but fits on a single page.)

## The Pak-Stanley labeling

## Remark

Mazin has shown that the Pak-Stanley labeling of the regions of the extended Shi arrangement is surjective. Together with Stanley's above result we have a self-contained proof of the fact that the Pak-Stanley labeling is a bijection between the regions of the regions of the extended Shi arrangement and the a-parking functions.

## Athanasiadis-Linusson diagrams

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## Definition

The regions of a contiguous, separated and integral deformation of the braid arrangement
$\left\{x_{i}-x_{j}=m: 1 \leq i<j<n, m \in[-\beta(j, i), \beta(i, j)]\right\}$ have Athanasiadis-Linusson diagrams if $\{\beta(i, j): i \neq j\}$ contains at most two consecutive nonnegative integers for each $j \in\{1,2, \ldots, n\}$. We set $\beta(j)=\min _{i \neq j} \beta(i, j)$ for all $j$.

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(2) For each $j$ satisfying $\beta(j)>0$ we also mark
$x_{j}+\beta(j), x_{j}+\beta(j)-1, \ldots, x_{j}+1$ on the reversed number line and we draw an arc connecting $x_{j}+k+1$ with $x_{j}+k$ for $k=0,1, \ldots, \beta(j)-1$. We label all of these points with $j$.

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(3) For each $\{i, j\} \subseteq\{1,2, \ldots, n\}$ we also draw an arc between $x_{i}$ and $x_{j}+\beta(j)$ if $\beta(i, j)=\beta(j)+1 x_{i}-x_{j}>\beta(i, j)$ holds.

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(9) We remove all nested arcs, that is, all arcs that contain another arc.

## Athanasiadis-Linusson diagrams

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## Athanasiadis-Linusson diagrams



Without 5 this is an example of Athanasiadis and Linusson in $\mathcal{A}_{3}^{1,2}$. For all $\{i, j\} \subset\{1,2,3,4\}$ we have $\beta(i, j)=2$ if $i<j$ and $\beta(i, j)=1$ if $i>j$. We add $\beta(i, 5)=\beta(5, i)=0$ for $i=1,2,4$, and we add $\beta(3,5)=1$ and $\beta(3,5)=0$.

## Athanasiadis-Linusson diagrams



For each $i \in\{1,2, \ldots, n\}$ we define $f(i)$ as the position of the leftmost element of the continuous component of $i$. We call the resulting $(f(1), f(2), \ldots, f(n))$ the $\beta$-parking function of the region.

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## Athanasiadis-Linusson trees

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(1) Replace the labels $j$ with $j_{1}, j_{2}, \ldots, j_{\beta(j)+1}$, numbered left to right, so that we can distinguish the copies.

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(2) The copies of the labels satisfying $f(j)=1$ become the children of the root 0 .
(3) We number the nodes in the tree level-by-level and in increasing order of the labels (breadth-first-search order).
(9) Once we inserted the copies of all labels $j$ satisfying $f(j)<i$, all copies of the labels $j$ satisfying $f(j)=i$ will be the children of the node whose number is $i$.

## Athanasiadis-Linusson trees

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## Definition

For a sequence $\underline{\beta} \in \mathbb{N}^{n}$ we define the $\underline{\beta}$-extended Shi arrangement as the hyperplane arrangement

$$
x_{i}-x_{j}=-\beta(j),-\beta(j)+1, \ldots, \beta(j)+1 \quad 1 \leq i<j \leq n \quad \text { in } V_{n-1}
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## Theorem

The number of regions in a $\underline{\beta}$-extended Shi arrangement $\mathcal{A}$ is

$$
r(\mathcal{A})=\left(\sum_{j=1}^{n}(\beta(j)+1)+1\right)^{n-1}
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The proof uses a colored variant of the Prüfer code algorithm.

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## a-Catalan arrangements

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The Athanasiadis-Linusson diagrams are very simple: they connect points with the same label only.

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$x_{\sigma(1)}>x_{\sigma(2)}>\cdots>x_{\sigma(n)}$, the parking trees are in bijection with the rooted incomplete $a$-ary trees on $(a-1) n+1$ vertices. Their number is the a-Catalan number $\frac{1}{(a-1) n+1}\binom{a n}{n}$. Multiplying it with $n$ ! we get

$$
r\left(\mathcal{A}_{n-1}^{a, a}\right)=a n(a n-1) \cdots((a-1) n+2)
$$

first found by Postnikov and Stanley.

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Fix a permutation $\pi$ and an a-Catalan path $\Lambda$.

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$$
w(\pi(i), \pi(j))= \begin{cases}\ell(\pi(j))-\ell(\pi(i)) & \text { if } \ell(\pi(j))-\ell(\pi(i)) \in[1-a, a-1] \\ -\infty & \text { if } \ell(\pi(j))-\ell(\pi(i))<1-a \\ a-1 & \text { if } \ell(\pi(j))-\ell(\pi(i))>a-1\end{cases}
$$

## A mysterious labeling

Fix a permutation $\pi$ and an a-Catalan path $\Lambda$.


## Lemma

The total order of gains $\sigma=\gamma \circ \pi$ is the order of the labels $\pi(1), \ldots, \pi(n)$ in increasing order of their levels, where $\pi(i)$ is listed before $\pi(j)$ if $\ell(\pi(i))=\ell(\pi(j))$ and $i<j$ hold.

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## Proposition

For the weighted digraph encoded by $(\pi, \Lambda)$ the gain function is the level function: we have $g(\sigma(i))=\ell(\sigma(i))$.

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The correspondence between the pairs $(\pi, \Lambda)$ and the valid weighted m-acyclic digraphs encoded by them is a bijection.

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We only prove injectivity and then we use the Postnikov-Stanley formula.

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## Proposition

A region of $\mathcal{A}_{n-1}^{a, a}$ is bounded if and only if the total order of gains $\sigma$ satisfies $w(\sigma(i), \sigma(i+1))<a-1$ for $1 \leq i \leq n-1$.

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The number of possible types of the trees of the gain function is a Catalan number.

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## Conjecture

For a fixed $n$ and a fixed tree of gain functions, the number of regions of $\mathcal{A}_{n-1}^{a, a}$ associated to it is a polynomial of a.

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The number of possible types of the trees of the gain function is a Catalan number.

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For a fixed $n$ and a fixed tree of gain functions, the number of regions of $\mathcal{A}_{n-1}^{a, a}$ associated to it is a polynomial of a.

This conjecture implies that the $n$-th a-Catalan number, considered as a polynomial of $a$, could be written as a sum of $C_{n}$ polynomials, where $C_{n}$ is the $n$-th Catalan number.

The general setup Sparse deformations Separated deformations

## Thank you!

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Labeling regions in deformations of graphical arrangements

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Labeling regions in deformations of graphical arrangements arXiv:2312.06513 [math.CO]

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