

Completeness Theorem:

Fact 1: $A \models \alpha$ iff $A \models \alpha$. If $A = \emptyset$, then $\models \alpha$ means that α is a tautology.

Propositional formula α is indecomposable iff α is a literal or negation of a literal

Example: $A, \sim A$

Sequence of formulas is **fundamental** if it contains a formula and its negation.

Example 1: $A, A \rightarrow (B \vee C), A \rightarrow C, B \vee A, \sim(A \rightarrow C), C \vee A$.

RS -strategy converts a formula α to a decomposition tree $T(\alpha)$ using 7 decomposition rules.

Disjunction decomposition rules

$$(\cup) \frac{\Gamma', (A \cup B), \Delta}{\Gamma', A, B, \Delta}, \quad (\neg \cup) \frac{\Gamma', \neg(A \cup B), \Delta}{\Gamma', \neg A, \Delta : \Gamma', \neg B, \Delta}$$

Conjunction decomposition rules

$$(\cap) \frac{\Gamma', (A \cap B), \Delta}{\Gamma', A, \Delta : \Gamma', B, \Delta}, \quad (\neg \cap) \frac{\Gamma', \neg(A \cap B), \Delta}{\Gamma', \neg A, \neg B, \Delta}$$

Implication decomposition rules

$$(\Rightarrow) \frac{\Gamma', (A \Rightarrow B), \Delta}{\Gamma', \neg A, B, \Delta}, \quad (\neg \Rightarrow) \frac{\Gamma', \neg(A \Rightarrow B), \Delta}{\Gamma', A, \Delta : \Gamma', \neg B, \Delta}$$

Negation decomposition rule

$$(\neg \neg) \frac{\Gamma', \neg \neg A, \Delta}{\Gamma', A, \Delta}$$

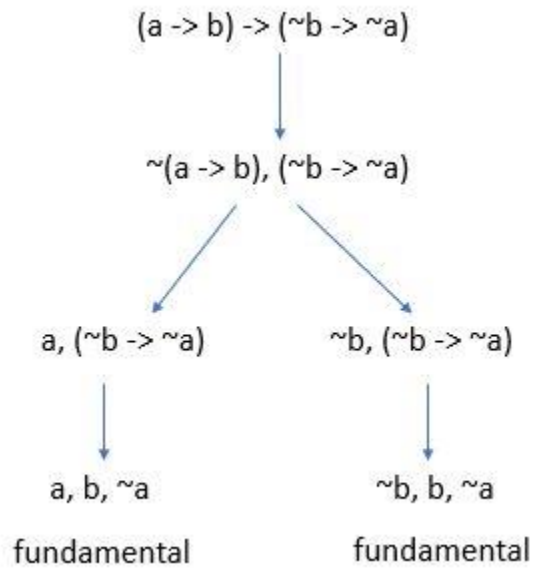
where Γ', Δ are sequences of formulas and Γ' is indecomposable (all formulas in Γ' are indecomposable)

Decomposition tree **$T(\alpha)$** , for a formula α is built from left to right!!!!

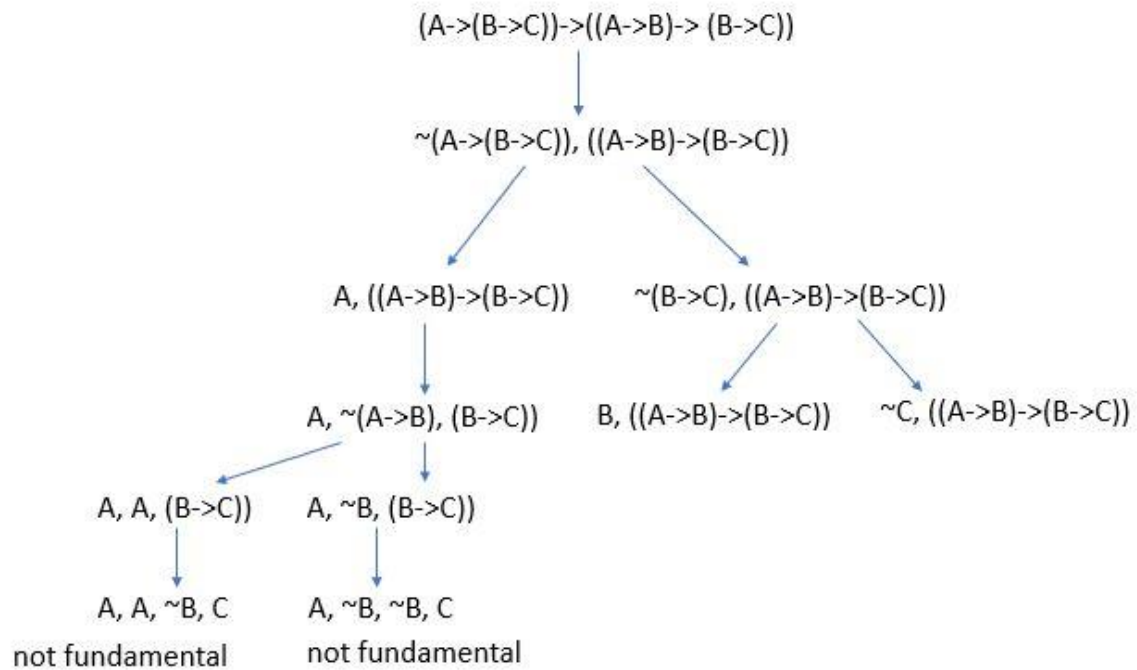
Theorem:

$\models \alpha$ iff every leaf in $T(\alpha)$ is fundamental.

Example 2: Decomposition tree $T((a \rightarrow b) \rightarrow (\neg b \rightarrow \neg a))$.



Example 3: Decomposition tree $T((A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (B \rightarrow C)))$.



Example 4:

Build decomposition tree for $(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$ using RS decomposition rules.

Remark: Use notation “;” when a node splits to two nodes.

$(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$\sim(A \rightarrow (B \rightarrow C)), ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$A, ((A \rightarrow B) \rightarrow (A \rightarrow C)); \sim(B \rightarrow C), ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$A, \sim(A \rightarrow B), (A \rightarrow C); \sim(B \rightarrow C), ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$A, A, (A \rightarrow C); A, \sim B, (A \rightarrow C); \sim(B \rightarrow C), ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$A, A, \sim A, C; A, \sim B, \sim A, C; \sim(B \rightarrow C), ((A \rightarrow B) \rightarrow (A \rightarrow C))$

fundamental; fundamental;

So, we continue with

$\sim(B \rightarrow C), ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$B, ((A \rightarrow B) \rightarrow (A \rightarrow C)); \sim C, ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$B, \sim(A \rightarrow B), (A \rightarrow C); \sim C, ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$B, A, (A \rightarrow C); B, \sim B, (A \rightarrow C); \sim C, ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$B, A, \sim A, C; B, \sim B, (A \rightarrow C); \sim C, ((A \rightarrow B) \rightarrow (A \rightarrow C))$

fundamental; fundamental;

So, we continue with

$\sim C, ((A \rightarrow B) \rightarrow (A \rightarrow C))$

$\sim C, \sim(A \rightarrow B), (A \rightarrow C)$

$\sim C, A, (A \rightarrow C); \sim C, \sim B, (A \rightarrow C)$

$\sim C, A, \sim A, C; \sim C, \sim B, \sim A, C$

fundamental; fundamental

So, all leaves of the decomposition tree are fundamental which means

$(A \rightarrow (B \rightarrow C)) \rightarrow ((A \rightarrow B) \rightarrow (A \rightarrow C))$ is a tautology.