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01001110010000110101001101010101

# Number Systems

4e435355

1,313,035,093

ECGR2181  
Lecture Notes 2

$1.525 \times 2^{29}$

*Reading:* Chapter 2

# How do we represent data in a computer?

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At the lowest level, a computer is an electronic machine.

- works by controlling the flow of electrons

Easy to recognize two conditions:

1. presence of a voltage – we'll call this state "1"
2. absence of a voltage – we'll call this state "0"

Could base state on *value* of voltage,  
but control and detection circuits more complex.

- compare turning on a light switch to  
measuring or regulating voltage

We'll see examples of these circuits in the next chapter.

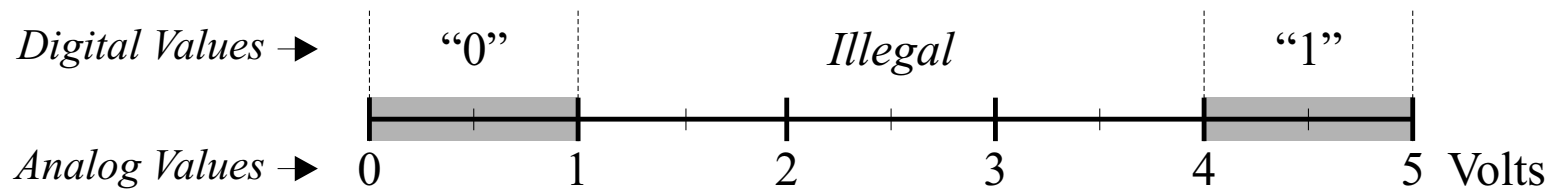
# Computer is a binary digital system.

Digital system:

- finite number of symbols

Binary (base two) system:

- has two states: 0 and 1



Basic unit of information is the *binary digit*, or **bit**.

Values with more than two states require multiple bits.

- A collection of two bits has four possible states:  
00, 01, 10, 11
- A collection of three bits has eight possible states:
- A collection of  $n$  bits has  $2^n$  possible states.

# What kinds of data do we need to represent?

- **Numbers** – signed, unsigned, integers, floating point, complex, rational, irrational, ...
- **Text** – characters, strings, ...
- **Images** – pixels, colors, shapes, ...
- **Sound**
- **Logical** – true, false
- **Instructions**
- ...

Data type:

- *representation and operations* within the computer

We'll start with numbers...

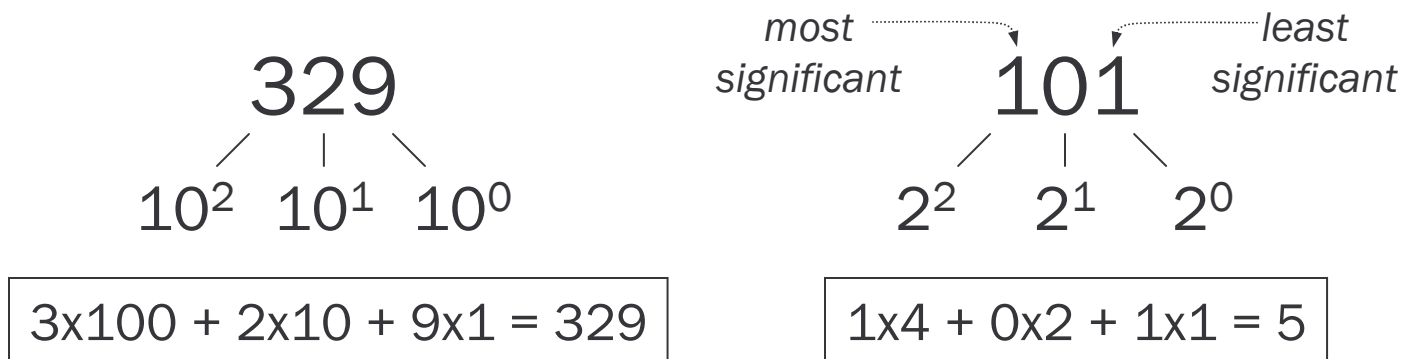
# Unsigned Integers

## Non-positional notation

- could represent a number (“5”) with a string of ones (“11111”)
- problems?

## Weighted positional notation

- like decimal numbers: “329”
- “3” is worth 300, because of its position, while “9” is only worth 9



## Unsigned Integers (cont.)

An  $n$ -bit unsigned integer represents  $2^n$  values:  
from 0 to  $2^n - 1$ .

$2^2$	$2^1$	$2^0$	Decimal
0	0	0	0
0	0	1	1
0	1	0	2
0	1	1	3
1	0	0	4
1	0	1	5
1	1	0	6
1	1	1	7

# Unsigned Integer Practice

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What is the decimal value of the unsigned binary number “1101”?

What is the binary representation of decimal 10?

# Unsigned Binary Arithmetic

Base-2 addition – just like base-10!

- add from right to left, propagating carry

$$\begin{array}{r} 10010 \\ + 1001 \\ \hline 11011 \end{array}$$
$$\begin{array}{r} 10010 \\ + 1011 \\ \hline 11101 \end{array}$$

*carry*

$$\begin{array}{r} 1111 \\ + 1 \\ \hline 10000 \end{array}$$



# Unsigned Binary Arithmetic Practice

---

Base-2 addition – just like base-10!

- add from right to left, propagating carry

$$\begin{array}{r} 10111 \\ + \underline{1111} \end{array}$$

$$\begin{array}{r} 10111 \\ - \underline{1111} \end{array}$$

Subtraction, multiplication, division,...

# Signed Integers

With  $n$  bits, we have  $2^n$  distinct values.

- assign about half to positive integers (1 through  $2^{n-1}$ ) and about half to negative ( $-2^{n-1}$  through  $-1$ )
- that leaves two values: one for 0, and one extra

## Positive integers

- just like unsigned – zero in most significant bit

$$00101 = 5$$

## Negative integers

- sign-magnitude – set top bit to show negative, bottom as in unsigned

$$10101 = -5$$

- one's complement – flip every bit to represent negative

$$11010 = -5$$

- in either case, most significant bit indicates sign:  
0=positive, 1=negative

# Is there a better way?

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## Problems with sign-magnitude and 1's complement

- two representations of zero (+0 and -0)
- arithmetic circuits are complex
  - How to add two sign-magnitude numbers?
    - e.g., try  $2 + (-3)$
  - How to add to one's complement numbers?
    - e.g., try  $4 + (-3)$

There is a great representation of negative numbers that uses the analogy of an automobile odometer.

# Odometer numbers

Consider an odometer of a car at a location on a street:

0	0	1	3
---	---	---	---

 Go 3 miles in reverse and it reads: 

0	0	1	0
---	---	---	---

- u Same as “subtracting 3” or “adding -3”
- u What happens when...

0	0	0	0
---	---	---	---

 Go 3 miles in reverse and it reads: 

9	9	9	7
---	---	---	---

- u As far as the odometer is concerned,  $9997 = -3$
- u Note that fixed-width binary is very similar to odometer numbers in its limitations
- u Can the same representation be used?
  - ◇  $0000_2 - 0001_2 = “1111_2”$  or -1
  - ◇  $0000_2 - 0011_2 = “1101_2”$  or -3
- u This is called “2’s complement”

# Two's Complement

*Two's complement* representation developed to make circuits easy for arithmetic.

- for each positive number (X), assign value to its negative (-X), such that  $X + (-X) = 0$  with “normal” addition, ignoring carry out

$$\begin{array}{r} 00101 \quad (5) \\ + \underline{11011} \quad (-5) \\ \hline 00000 \quad (0) \end{array}$$

$$\begin{array}{r} 01001 \quad (9) \\ + \underline{\quad\quad\quad} \quad (-9) \\ \hline 00000 \quad (0) \end{array}$$

# Two's Complement Representation

If number is positive or zero,

- normal binary representation, zeroes in upper bit(s)

If number is negative,

- start with positive number
- flip every bit (i.e., take the one's complement)
- then add one

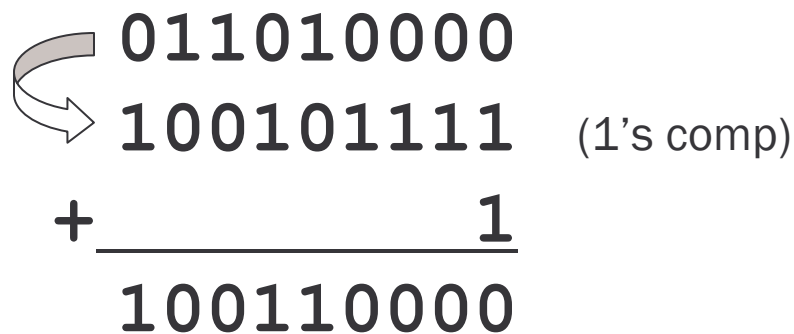
$$\begin{array}{r} \curvearrowright \text{00101} \quad (5) \\ \curvearrowleft \text{11010} \quad (1\text{'s comp}) \\ + \quad \underline{\quad 1} \\ \text{11011} \quad (-5) \end{array}$$

$$\begin{array}{r} \curvearrowright \text{01001} \quad (9) \\ \curvearrowleft \quad \quad \quad (1\text{'s comp}) \\ + \quad \underline{\quad 1} \\ \quad \quad \quad (-9) \end{array}$$

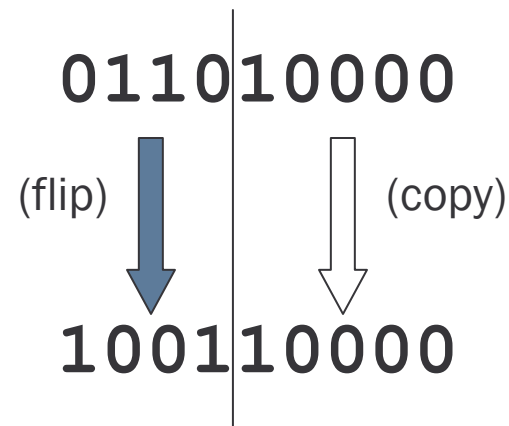
# Two's Complement Shortcut

To take the two's complement of a number:

- copy bits from right to left until (and including) the first “1”
- flip remaining bits to the left



011010000  
100101111 (1's comp)  
+ 1  
-----  
100110000



# Two's Complement Signed Integers

MS bit is sign bit – it has weight  $-2^{n-1}$ .

Range of an n-bit number:  $-2^{n-1}$  through  $2^{n-1} - 1$ .

- The most negative number ( $-2^{n-1}$ ) has no positive counterpart.

$-2^3$	$2^2$	$2^1$	$2^0$		$-2^3$	$2^2$	$2^1$	$2^0$	
0	0	0	0	0	1	0	0	0	-8
0	0	0	1	1	1	0	0	1	-7
0	0	1	0	2	1	0	1	0	-6
0	0	1	1	3	1	0	1	1	-5
0	1	0	0	4	1	1	0	0	-4
0	1	0	1	5	1	1	0	1	-3
0	1	1	0	6	1	1	1	0	-2
0	1	1	1	7	1	1	1	1	-1



# Two's Complement Practice

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Show the two's complement representation of the decimal number -6.

# Converting Binary (2's C) to Decimal

1. If leading bit is one, take two's complement to get a positive number.
2. Add powers of 2 that have "1" in the corresponding bit positions.
3. If original number was negative, add a minus sign.

$$\begin{aligned} X &= 01101000_{\text{two}} \\ &= 2^6 + 2^5 + 2^3 = 64 + 32 + 8 \\ &= 104_{\text{ten}} \end{aligned}$$

$n$	$2^n$
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
(K) 10	1024
(M) 20	1048576

*Assuming 8-bit 2's complement numbers.*

# More Examples

$$\begin{aligned}
 X &= 00100111_{\text{two}} \\
 &= 2^5 + 2^2 + 2^1 + 2^0 = 32 + 4 + 2 + 1 \\
 &= 39_{\text{ten}}
 \end{aligned}$$

$$\begin{aligned}
 X &= 11100110_{\text{two}} \\
 -X &= 00011010 \\
 &= 2^4 + 2^3 + 2^1 = 16 + 8 + 2 \\
 &= 26_{\text{ten}} \\
 X &= -26_{\text{ten}}
 \end{aligned}$$

$n$	$2^n$
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
(K) 10	1024
(M) 20	1048576

Assuming 8-bit 2's complement numbers.

# Converting Binary (2's C) to Decimal Practice

Convert binary 00011111 to decimal:

$n$	$2^n$
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
(K) 10	1024
(M) 20	1048576

# Converting Decimal to Binary (2's C)

## First Method: *Division*

1. Divide by two – remainder is least significant bit.
2. Keep dividing by two until answer is zero, writing remainders from right to left.
3. Append a zero as the MS bit;  
if original number negative, take two's complement.

$$X = 104_{\text{ten}}$$

$$104/2 = 52 \text{ r}0 \text{ bit } 0$$

$$52/2 = 26 \text{ r}0 \text{ bit } 1$$

$$26/2 = 13 \text{ r}0 \text{ bit } 2$$

$$13/2 = 6 \text{ r}1 \text{ bit } 3$$

$$6/2 = 3 \text{ r}0 \text{ bit } 4$$

$$3/2 = 1 \text{ r}1 \text{ bit } 5$$

$$X = 01101000_{\text{two}}$$

$$1/2 = 0 \text{ r}1 \text{ bit } 6$$

# Converting Decimal to Binary (2's C)

## Second Method: *Subtract Powers of Two*

1. Change to positive decimal number.
2. Subtract largest power of two less than or equal to number.
3. Put a one in the corresponding bit position.
4. Keep subtracting until result is zero.
5. Append a zero as MS bit; if original was negative, take two's complement.

$n$	$2^n$
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
(K) 10	1024
(M) 20	1048576

$$\begin{array}{rcl}
 X = 104_{\text{ten}} & 104 - 64 = 40 & \text{bit 6} \\
 & 40 - 32 = 8 & \text{bit 5} \\
 & 8 - 8 = 0 & \text{bit 3}
 \end{array}$$

$$X = 01101000_{\text{two}}$$

# Converting Decimal to Binary Practice

Convert decimal 270 to binary using both methods described above:

$n$	$2^n$
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
(K) 10	1024
(M) 20	1048576

# More Converting Decimal to Binary Practice

Convert decimal 255 to binary using both methods described above:

$n$	$2^n$
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
(K) 10	1024
(M) 20	1048576



# Operations: Arithmetic and Logical

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Recall: a data type includes *representation* and *operations*.

We now have a good representation for signed integers, so let's look at some ***arithmetic*** operations:

- Addition
- Subtraction
- Sign Extension

We'll also look overflow conditions for addition.

Multiplication, division, etc., can be built from these basic operations.

***Logical*** operations are also useful:

- AND
- OR
- NOT

# Addition

As we've discussed, 2's comp. addition is just binary addition.

- assume all integers have the same number of bits
- ignore carry out
- for now, assume that sum fits in n-bit 2's comp. representation

$$\begin{array}{r} 01101000 \quad (104) \\ + \underline{11110000} \quad (-16) \\ \hline 01011000 \quad (88) \end{array} \qquad \begin{array}{r} 11110110 \quad (-10) \\ + \underline{\hspace{1cm}} \quad (-9) \\ \hline \hspace{1cm} \quad (-19) \end{array}$$

*Assuming 8-bit 2's complement numbers.*

# Subtraction

Negate subtrahend (2nd no.) and add.

- assume all integers have the same number of bits
- ignore carry out
- for now, assume that difference fits in n-bit 2's comp. representation

$$\begin{array}{r} 01101000 \text{ (104)} \\ - 00010000 \text{ (16)} \\ \hline \end{array} \quad \begin{array}{r} 11110110 \text{ (-10)} \\ - \phantom{00000000} \text{ (-9)} \\ \hline \end{array}$$

is just

$$\begin{array}{r} 01101000 \text{ (104)} \\ + 11110000 \text{ (-16)} \\ \hline 01011000 \text{ (88)} \end{array} \quad \begin{array}{r} 11110110 \text{ (-10)} \\ + \phantom{00000000} \text{ (9)} \\ \hline \phantom{01011000} \text{ (-1)} \end{array}$$

*Assuming 8-bit 2's complement numbers.*

# Practice

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Perform the Two's Complement operation to the following decimal numbers: - 56 - 14

# Sign Extension

To add two numbers, we must represent them with the same number of bits.

If we just pad with zeroes on the left:

<u>4-bit</u>	<u>8-bit</u>
0100 (4)	00000100 (still 4)
1100 (-4)	00001100 (12, not -4)

Instead, replicate the most significant bit -- the sign bit:

<u>4-bit</u>	<u>8-bit</u>
0100 (4)	00000100 (still 4)
1100 (-4)	11111100 (still -4)

# Overflow

If operands are too big,  
then sum cannot be represented as an  $n$ -bit 2's comp  
number.

<b>01000</b>	(8)	<b>11000</b>	(-8)
<b>+ 01001</b>	(9)	<b>+ 10111</b>	(-9)
<b>10001</b>	(-15)	<b>01111</b>	(+15)

We have overflow if:

- signs of both operands are the same, and
- sign of sum is different.

Another test -- easy for hardware:

- carry into MS bit does not equal carry out

# Logical Operations

## Operations on logical TRUE or FALSE

- two states -- takes one bit to represent: TRUE=1, FALSE=0

## View $n$ -bit number as a collection of $n$ logical values

- operation applied to each bit independently

A	B	A <b>AND</b> B	A	B	A <b>OR</b> B	A	<b>NOT</b> A
0	0	0	0	0	0	0	1
0	1	0	0	1	1	1	0
1	0	0	1	0	1		
1	1	1	1	1	1		

# Examples of Logical Operations

## AND

- useful for clearing bits
  - AND with zero = 0
  - AND with one = no change

$$\begin{array}{r} \phantom{\text{AND}} \quad 11000101 \\ \text{AND} \quad \underline{00001111} \\ \phantom{\text{AND}} \quad 00000101 \end{array}$$

## OR

- useful for setting bits
  - OR with zero = no change
  - OR with one = 1

$$\begin{array}{r} \phantom{\text{OR}} \quad 11000101 \\ \text{OR} \quad \underline{00001111} \\ \phantom{\text{OR}} \quad 11001111 \end{array}$$

## NOT

- unary operation -- one argument
- flips every bit

$$\begin{array}{r} \text{NOT} \quad \underline{11000101} \\ \phantom{\text{NOT}} \quad 00111010 \end{array}$$



# Practice of Logical Operations

## AND

- useful for clearing bits
  - AND with zero = 0
  - AND with one = no change

                  10101010  
**AND**           11100011

## OR

- useful for setting bits
  - OR with zero = no change
  - OR with one = 1

                  11001100  
**OR**           00001111

## NOT

- unary operation -- one argument
- flips every bit

**NOT**           10010110

# Hexadecimal Notation

It is often convenient to write binary (base-2) numbers as hexadecimal (base-16) numbers instead.

- fewer digits -- four bits per hex digit
- less error prone -- easy to corrupt long string of 1's and 0's

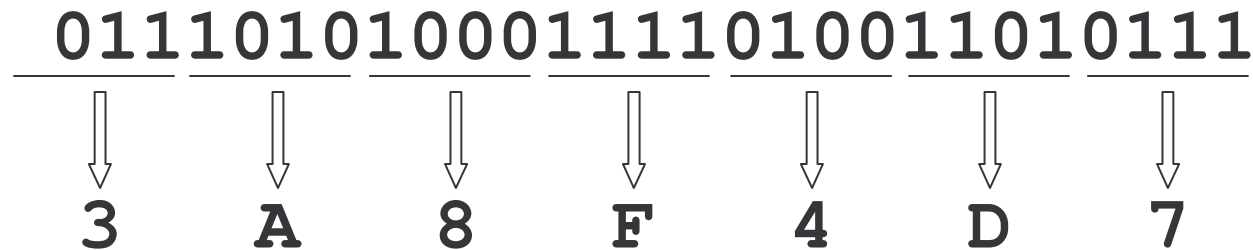
Binary	Hex	Decimal	Binary	Hex	Decimal
0000	0	0	1000	8	8
0001	1	1	1001	9	9
0010	2	2	1010	A	10
0011	3	3	1011	B	11
0100	4	4	1100	C	12
0101	5	5	1101	D	13
0110	6	6	1110	E	14
0111	7	7	1111	F	15

Memorize this table!!!!

# Converting from Binary to Hexadecimal

Every four bits is a hex digit.

- start grouping from right-hand side

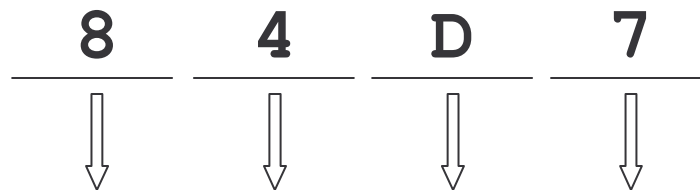


*This is not a new machine representation,  
just a convenient way to write the number.*

# Converting from Hexadecimal to Decimal

Every hex digit position has a base value

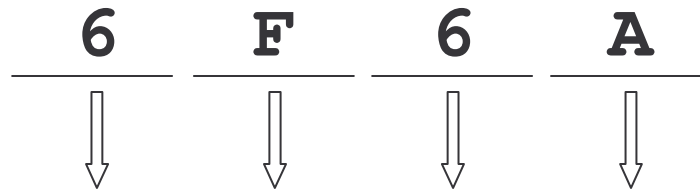
- multiply the value at the position by the base value



$$\begin{aligned} & 8 \times 16^3 + 4 \times 16^2 + 13 \times 16^1 + 7 \times 16^0 = \\ & 8 \times 4096 + 4 \times 256 + 13 \times 16 + 7 \times 1 = \\ & \quad 32768 + 1024 + 208 + 7 = \\ & \qquad \qquad \qquad 34007 \end{aligned}$$

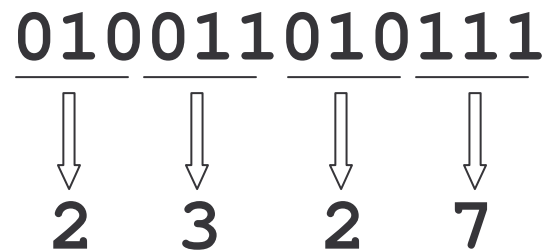
# Practice Converting from Hex to Decimal

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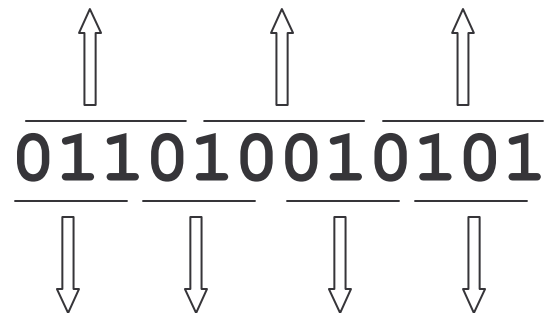


# Octal

Octal is simply “base 8” number representation



Octal and  
Hex  
Practice



# Multiplication by example

Consider  $2_{10} \times 3_{10}$

multiplicand	0010	
multiplier	0011	
	<hr/>	
	0010	copy of multiplicand
	0010	copy of multiplicand
	0000	zero
	0000	zero
	<hr/>	
	0000110	

# Text: ASCII Characters

ASCII: Maps 128 characters to 7-bit code.

- both printable and non-printable (ESC, DEL, ...) characters

00	nul	10	dle	20	sp	30	0	40	@	50	P	60	`	70	p
01	soh	11	dc1	21	!	31	1	41	A	51	Q	61	a	71	q
02	stx	12	dc2	22	"	32	2	42	B	52	R	62	b	72	r
03	etx	13	dc3	23	#	33	3	43	C	53	S	63	c	73	s
04	eot	14	dc4	24	\$	34	4	44	D	54	T	64	d	74	t
05	enq	15	nak	25	%	35	5	45	E	55	U	65	e	75	u
06	ack	16	syn	26	&	36	6	46	F	56	V	66	f	76	v
07	bel	17	etb	27	'	37	7	47	G	57	W	67	g	77	w
08	bs	18	can	28	(	38	8	48	H	58	X	68	h	78	x
09	ht	19	em	29	)	39	9	49	I	59	Y	69	i	79	y
0a	nl	1a	sub	2a	*	3a	:	4a	J	5a	Z	6a	j	7a	z
0b	vt	1b	esc	2b	+	3b	;	4b	K	5b	[	6b	k	7b	{
0c	np	1c	fs	2c	,	3c	<	4c	L	5c	\	6c	l	7c	
0d	cr	1d	gs	2d	-	3d	=	4d	M	5d	]	6d	m	7d	}
0e	so	1e	rs	2e	.	3e	>	4e	N	5e	^	6e	n	7e	~
0f	si	1f	us	2f	/	3f	?	4f	O	5f	_	6f	o	7f	del



# Data Communications

There was no standard for networks in the early days and as a result it was difficult for networks to communicate with each other.

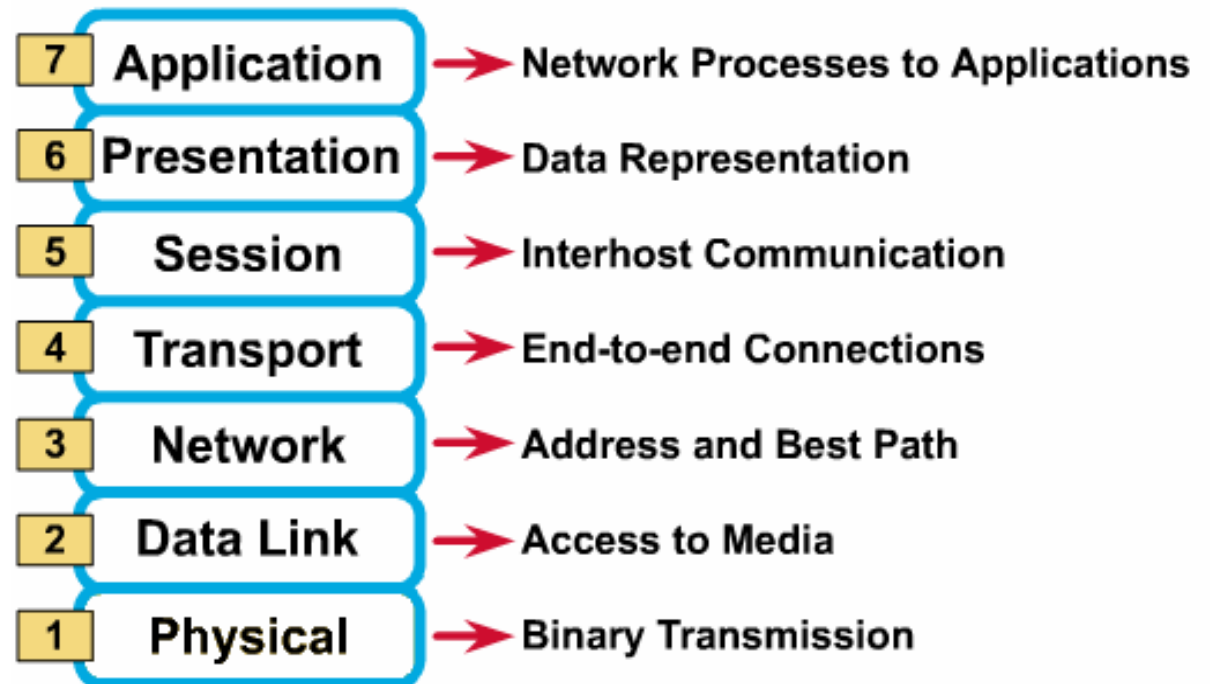
The International Organization for Standardization (ISO) recognized this and in 1984 introduced the Open Systems Interconnection (OSI) reference model.

The OSI reference model organizes network functions into seven numbered layers.

Each layer provides a service to the layer above it in the protocol specification and communicates with the same layer's software or hardware on other computers.

Layers 5-7 are concerned with services for the applications.

Layers 1-4 are concerned with the flow of data from end to end through the network



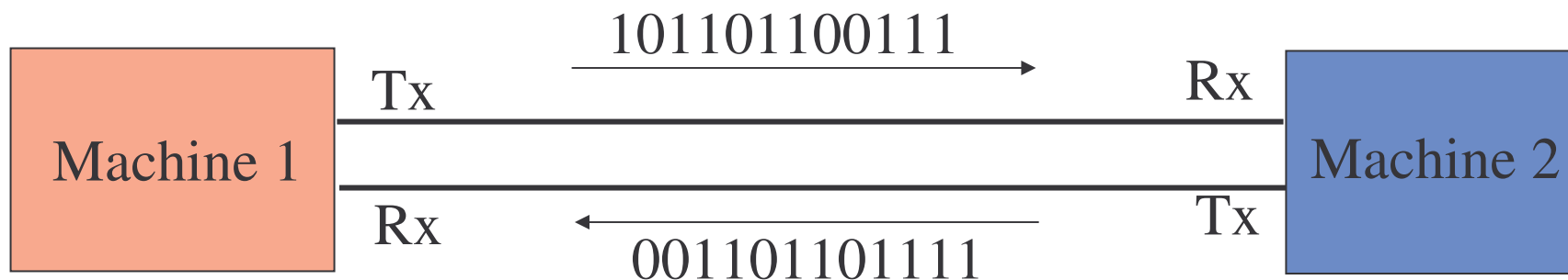
# Physical Layer (1) – Serial Communications

The basic premise of serial communications is that one or two wires are used to transmit digital data.

- Of course, ground reference is also needed (extra wire)

Can be one way or two way, usually two way, hence two communications wires.

Often other wires are used for other aspects of the communications (ground, “clear-to-send”, “data terminal ready”, etc).



# Serial Communication Basics

Send one bit of the message at a time

Message fields

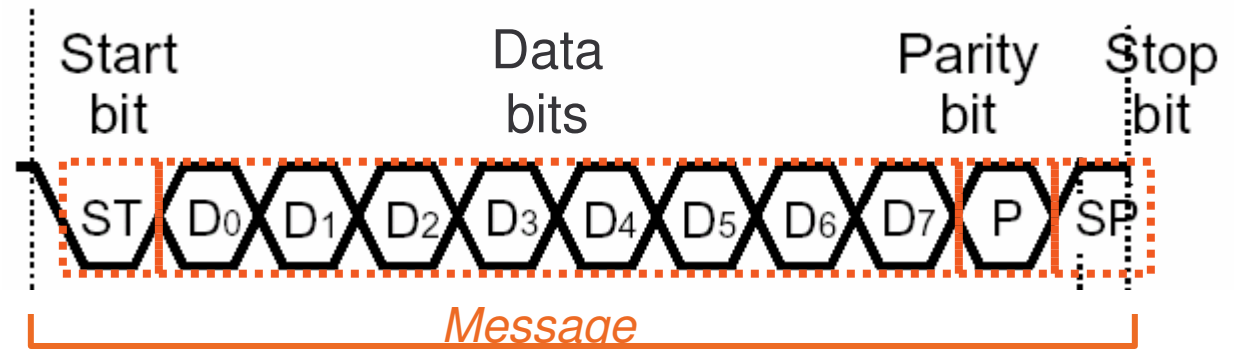
- Start bit (one bit)
- Data (LSB first or MSB, and size – 7, 8, 9 bits)
- Optional parity bit is used to make total number of ones in data even or odd
- Stop bit (one or two bits)

All devices on network or link must use same communications parameters

- The speed of communication must be the same as well (300, 600, 1200, 2400, 9600, 14400, 19200, etc.)

More sophisticated network protocols have more information in each message

- Medium access control – when multiple nodes are on bus, they must arbitrate for permission to transmit
- Addressing information – for which node is this message intended?
- Larger data payload
- Stronger error detection or error correction information
- Request for immediate response (“in-frame”)



# Serial Communication Basics

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RS232: rules on connector, signals/pins, voltage levels, handshaking, etc.

[RS232: Fulfilling All Your Communication Needs](#), Robert Ashby

[Quick Reference for RS485, RS422, RS232 and RS423](#)

Not so quick reference:

[The RS232 Standard: A Tutorial with Signal Names and Definitions](#),

Christopher E. Strangio

Bit vs Baud rates:

<http://www.totse.com/en/technology/telecommunications/bits.html>