

10/15/2009 Outline

- Friday 10/16/2009 Recitation starts at 8 AM for Tutorial on Tools and SW
- Exam 1 handed back on Friday
- HW 6 Due today
- HW 7 Due ~~10/22 Thursday~~ Friday 10/23

Chapter 4 - Digital Design

Adder :

$$\begin{array}{r} 01 \\ + 10 \\ \hline 11 \end{array}$$

$$\begin{array}{r} 01 \\ + 11 \\ \hline 100 \end{array}$$

2 bit adder $2^{(2+2)} = 16$ rows

8 bit adder $2^{(8+8)} = 65,536$ rows

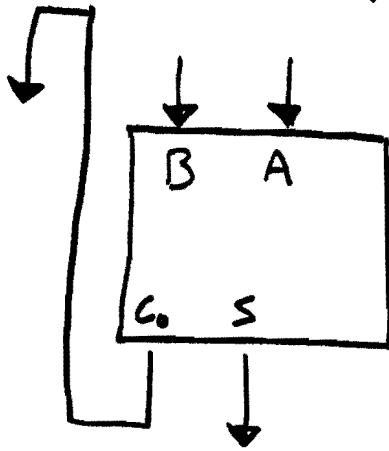
16 bit adder $2^{(16+16)} = 4$ Billion

Add 1 column at a time (like we humans do)

$$\begin{array}{r} 1111 \\ + 0110 \\ \hline 0111 \end{array} \quad \text{etc}$$

~~Exercise~~

Half Adder: (1-bit Adder
Sum and Carry out)

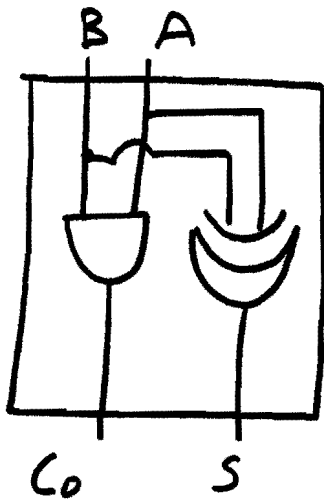


Capture the function

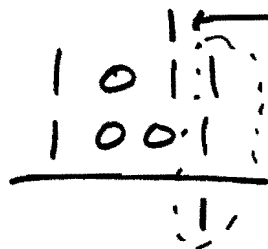
A	B	Co	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

$$C_0 = A \cdot B$$

$$S = A \text{ XOR } B = (A'B + AB')$$

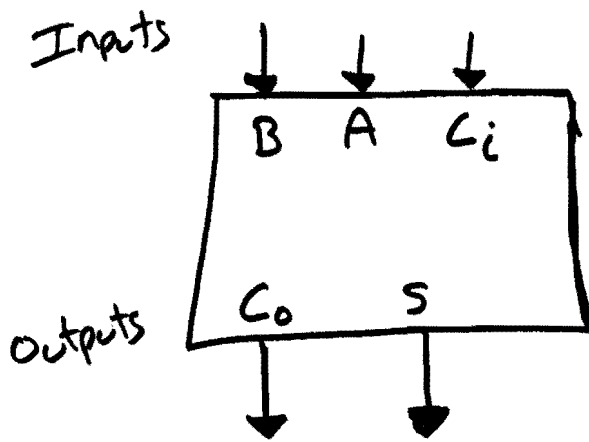


Limitation: multiple bit Half Adder does not allow us to consider Carry out



Full Adder

Adds 3 bits



Inputs			Outputs	
A	B	C_i	C_o	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

$$C_o = \bar{a}bc + a\bar{b}c + ab\bar{c} + abc$$

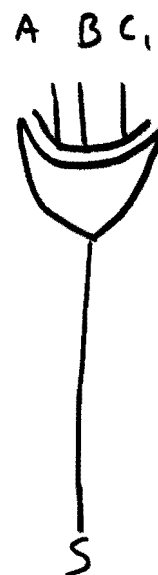
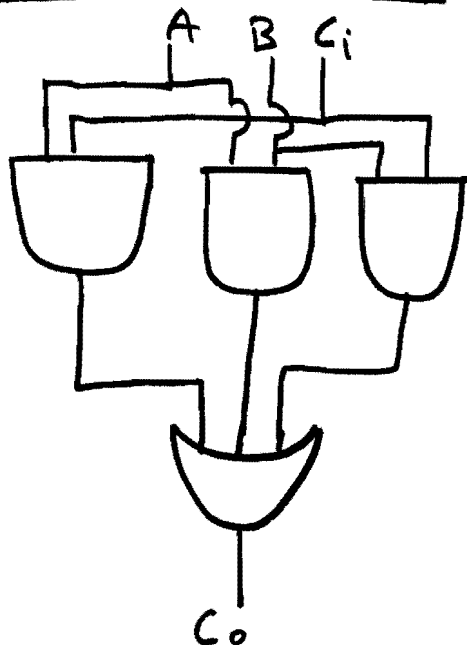
$$C_o = bc(\bar{a} + a) + a\bar{b}c + ab\bar{c}$$

$$C_o = bc + a(\bar{b}c + b\bar{c})$$

$$bc + a(c + b)$$

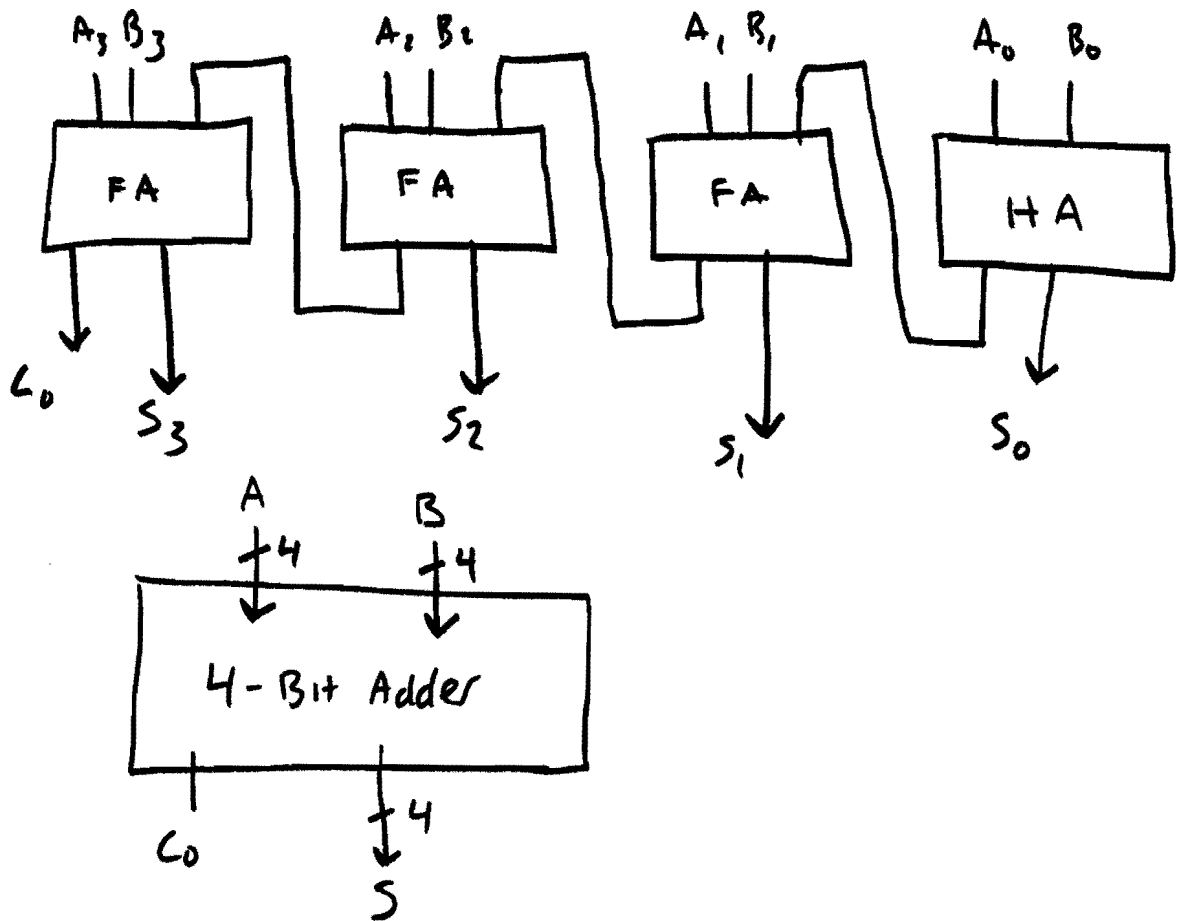
$$C_o = bc + ac + ab$$

$$S = a \text{ xor } b \text{ xor } C_i$$



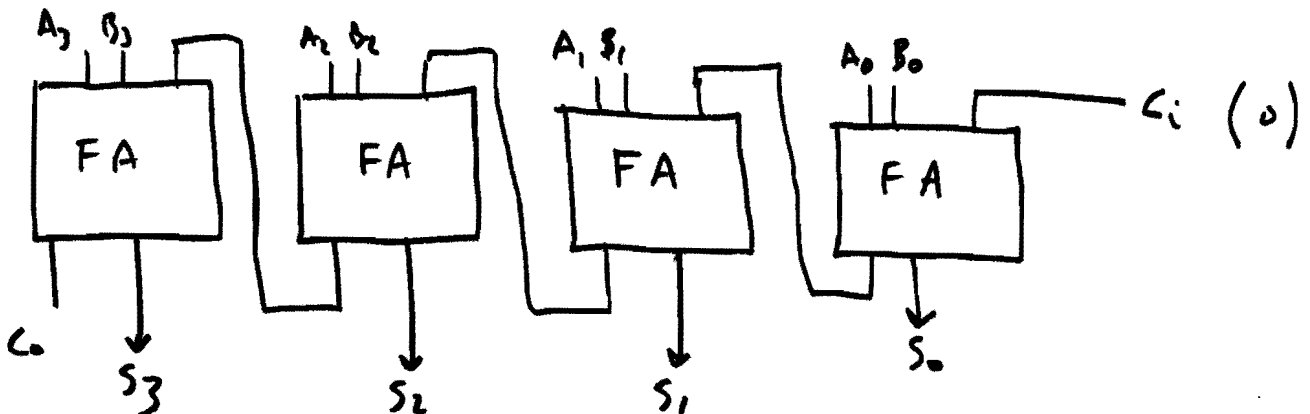
Carry-Ripple Adder v1.0

Multiple Bit Adder



Carry-Ripple Adder v2.0

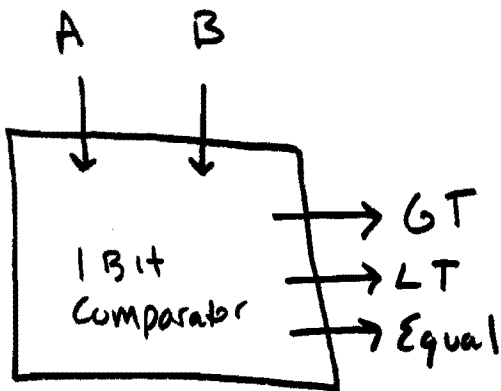
4 - Full Adders



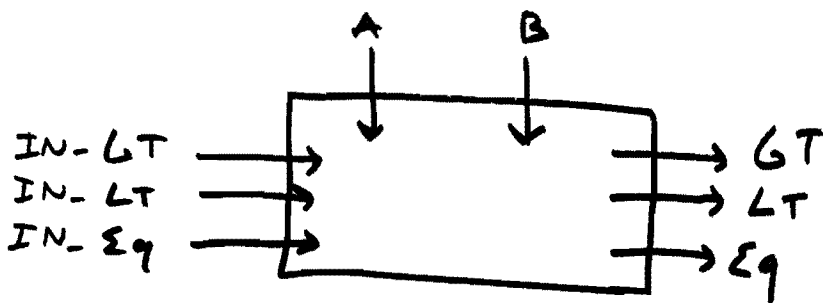
Magnitude Comparator

$A = 1011$
 Consider: $B = 1001$
 equal
 unequal
 $A > B$
 $1 > 0$

$A = 1011$
 $B = 1011$
 $A = B$



How to Chain multiple together?



Initial input is ~~A=B~~ $Eq = 1$

Comparator Greater Than? $A > B$

$$\text{OUT-GT} = \text{IN-GT} + (\text{IN-EQ} \cdot \underbrace{A \cdot B'}_{A > B})$$

$A = 1, B = 0$

Comparator Less Than? $A < B$

$$\text{OUT-LT} = \text{IN-LT} + (\text{IN-EQ} \cdot \underbrace{A' \cdot B}_{A < B})$$

$A = 0, B = 1$

Already know how to do an equality comparator ✓

$$\text{OUT-EQ} = \text{IN-EQ} \cdot (A \text{ XOR } B)$$

Worst Case Propagation Delay?

$$\text{AND, OR, NOT} = 1 T_p$$

$$\text{XOR} = 2 T_p$$

1 cell (1 bit comparator): Equality uses 1 XOR
1 AND

$$2 T_p + 1 T_p = 3 T_p$$

$$3 T_p + 2 \cdot (1 T_p) + 2 T_p$$

1st cell

middle cells

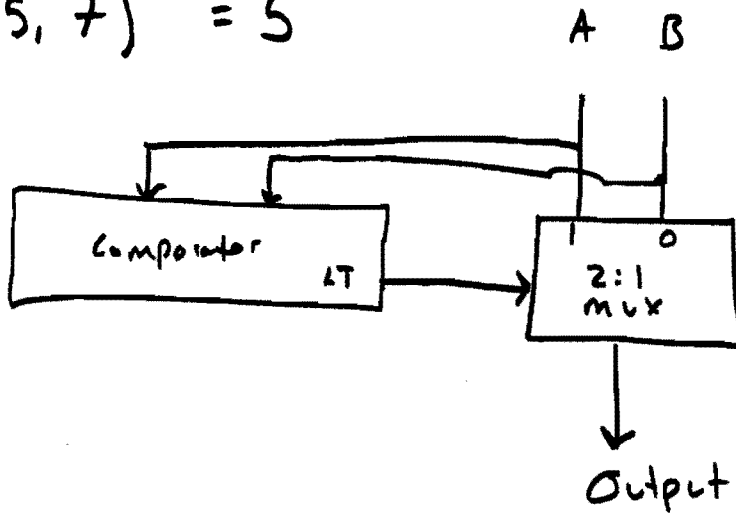
last cell

> or <

Minimum of Two Numbers

$$\text{min}(A, B)$$

$$\text{min}(5, 7) = 5$$



$$\underline{A < B}$$

if yes then
mux select bit = 1
so output = A
else
output = B

Multiplier

Slide 29

$$\begin{array}{r} 6 \\ \times 3 \\ \hline 18 \end{array}$$

$$\begin{array}{r} 0110 \\ \times 0011 \\ \hline 0110 \\ 0110 \\ 0000 \\ 0000 \\ \hline 0010010 \end{array}$$

AND

Then Add together

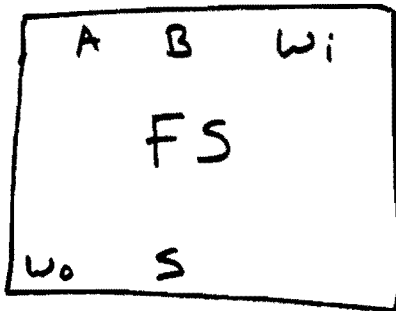
$$10010_2 = 16 + 2 = \boxed{18}$$

A_3	A_2	A_1	A_0
B_3	B_2	B_1	B_0
<hr/>			
$A_3 \cdot B_0$	$A_2 \cdot B_0$	$A_1 \cdot B_0$	$A_0 \cdot B_0$
$A_3 \cdot B_1$	$A_2 \cdot B_1$	$A_1 \cdot B_1$	$A_0 \cdot B_1$
\vdots			\vdots

AND GATES Feeding into Adders

Subtractor

Not using 2's complement



w_1 is borrowed from Right
 w_0 borrowed on left

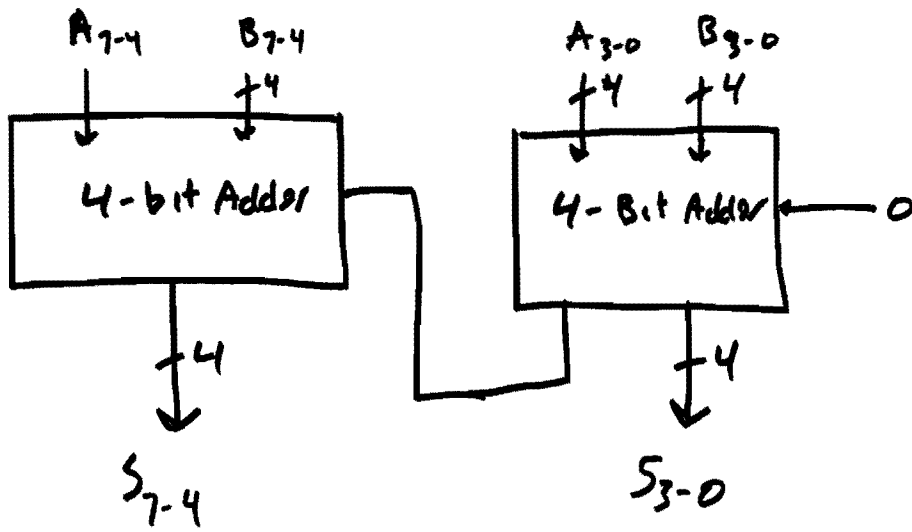
A	B	w_i	w_o	S
0	0	0	0	0
0	0	1	0	0
0	1	0	1	0
0	1	1	0	0
1	0	0	0	1
1	0	1	0	1
1	1	0	0	0
1	1	1	0	0

2's complement Adder

$A + B$ is Adder

$A - B$ is $A + (-B)$

Cascading Adders



Slides from Ch. 4 (up to 17)

Comparators :

Equality Comparator

$A = B$ iff $A=0 + B=0$
or
 $A=1 + B=1$

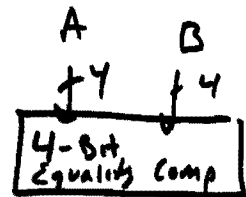
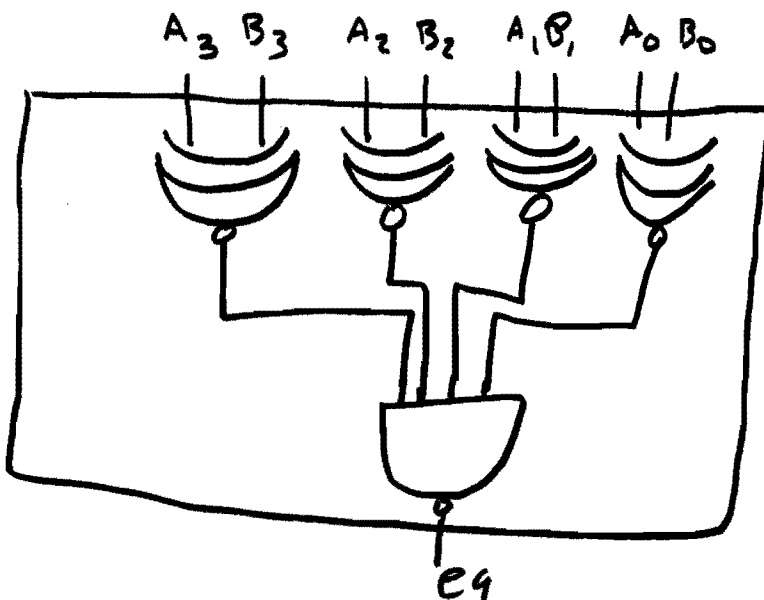
What Gate does this?

A	B	C
0	0	1
0	1	0
1	0	0
1	1	1

XNOR

$$A \oplus B = C$$

4 Bit Equality Comparator ?



$e_q = 1$ = Yes
 $e_q = 0$ = No